

Draw out and label the four parts of memory and mark how much is allocated in each section (except for code, whose size you don't know) at the point in execution in which the program uses the most memory. Assume a 32 bit system. (Hint: Think about which registers the function needs to preserve)

*0 bytes for preprocessor defines*  
*4 bytes int, 4 bytes char\* in static*  
*31 bytes for string in heap*  
*foo loops through 32 times (4bytes \* 32)*  
*Doesn't need to save \$a0 or \$v0 b/c they're not needed after*  
*Total = 31+8+128= 167 bytes*

### Number Representation:

[illegible]

**Ans:  $a > c > e > d > f > b$**

## MIPS:

Be a compiler. Translate the following C into MIPS TAL:

```
int dot_product (pixel* pix1, pixel* pix2)
{
    int temp;
    pix1->r *= pix2->r;
    pix1->g *= pix2->g;
    pix1->b *= pix2->b;

    temp = (pix1->r + pix1->g + pix1->b) / 3;

    if (temp <= 255)
    {
        pix1->moy = temp;
        return 0;
    }
    return 1;
}
```

Assume that the following data type has been declared and translated for you:

```
typedef struct colour
{
    short r;
    short g;
    short b;
    char moy;
} pixel;
```

## Ans:

```
dot_product: addi $v0, $0, 1
              lhu $t0, 0($a0) #r
              lhu $t1, 0($a1) #r2
              mult $t0, $t1
              mflo $t7
              sh $t7, 0($a0)

              lhu $t0, 2($a0) #g
              lhu $t1, 2($a1) #g2
              mult $t0, $t1
              mflo $t8
              sh $t8, 2($a0)

              lhu $t0, 4($a0) #b
              lhu $t1, 4($a1) #b2
              mult $t0, $t1
              mflo $t9
              sh $t9, 4($a0)

              addu $t7, $t7, $t8
              addu $t7, $t7, $t9
              addi $t0, $0, 3
              div $t7, $t0
              mflo $t7
              sltiu $t0, $t7, 256
              beq $t0, $0, end
              sb $t7, 6($a0)
              add $v0, $0, $0
end:          jr $ra
```

## CALL:

\* T/F: If a program (a single .c file with a main() statement) does not rely on external libraries or files, it does not need to undergo the linker step. As in, it can be loaded and executed immediately after the assembler step. *{F – object file cannot be loaded even though it's machine code; certain things have not been added, e.g. table of local symbols, additional bits the loader needs have not been added}*

### The following 3 questions apply to C programs on computers without software emulation

\* The resulting code is portable between different operating systems (e.g. Solaris and Windows) (circle all that apply)...

before compilation  
after compilation but before assembling  
after assembling but before linking  
after linking but before loading

***Ans: All of the above. Code is portable all the way up to the binary level (look at Wine for example)***

\* The resulting code is portable between different hardware ISAs (e.g. x86, MIPS, SPARC, and PowerPC) (circle all that apply)...

before compilation  
after compilation but before assembling  
after assembling but before linking  
after linking but before loading

***Ans: Only compilation. Once compiled to machine code, it cannot be moved to another ISA. An exception is Apple's Rosetta, but that is a form of "software emulation."***

\* T/F: If we compile & assemble a program that uses dynamic libraries on Solaris, can we move that object file to a MS Windows computer and link with the Windows versions of those libraries? *{F – ELF format vs MS proprietary PE format}*